



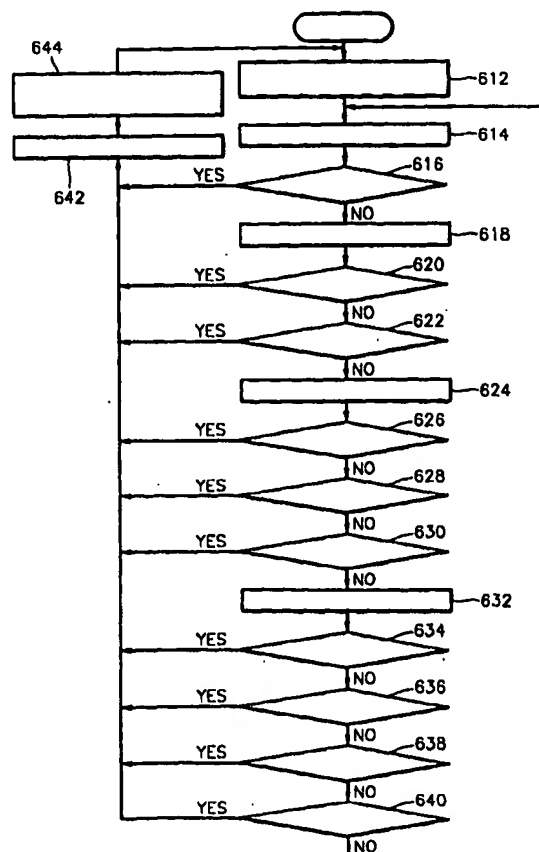
## INTERNATIONAL APPLICATION PUBLISHED UNDER THE PATENT COOPERATION TREATY (PCT)

(51) International Patent Classification <sup>6</sup> : <b>H03M 13/00</b>		<b>A1</b>	(11) International Publication Number: <b>WO 98/31106</b>
			(43) International Publication Date: 16 July 1998 (16.07.98)
(21) International Application Number: PCT/KR98/00004		(81) Designated States: AU, BR, CA, CN, DE, ES, GB, JP, RU, SG, European patent (AT, BE, CH, DE, DK, ES, FI, FR, GB, GR, IE, IT, LU, MC, NL, PT, SE).	
(22) International Filing Date: 12 January 1998 (12.01.98)			
(30) Priority Data: 08/782,174 14 January 1997 (14.01.97) US		Published With international search report.	
(71) Applicants: SAMSUNG ELECTRONICS CO., LTD. [KR/KR]; 416, Maetan-dong, Paldal-gu, Suwon-City, Kyungki-do 442-373 (KR). REGENTS OF THE UNIVERSITY OF CALIFORNIA [US/US]; 22nd floor, The 300 Lakeside Drive, Oakland, CA 94612-3550 (US).			
(72) Inventors: PARK, Dong, Seek; 1424-30, Sankyeok 4-dong, Buk-gu, Daegu-City 702-014 (KR). VILLASENOR, John; 56-125FF Engineering Building IV, Box 951594, Los Angeles, CA 90095 (US). CHEN, Feng; 56-125FF Engineering Building IV, Box 951594, Los Angeles, CA 90095 (US). DOWLING, Drendan; 56-125FF Engineering Building IV, Box 951594, Los Angeles, CA 990095 (US). LULTERELL, Max; 56-125FF Engineering Building IV, Box 951594, Los Angeles, CA 90095 (US).			
(74) Agent: LEE, Young, Pil; The Cheonghwa Building, 1571-18, Seocho-dong, Seocho-gu, Seoul 137-073 (KR).			

(54) Title: ERROR PREVENTION METHOD FOR MULTIMEDIA

## (57) Abstract

An error prevention method for multimedia improves data recovery and channel throughput in channels which cause a random error and a burst error by using a rate compatible punctured convolutional code (RCPC) and an automatic retransmission on request (ARQ). In a process of decoding a plurality of packets of given information, the error prevention method includes the steps of a) decoding one of the plurality of packets, b) decoding another packet when an error occurs during the decoding in step a), c) decoding a combination of the packets from steps a) and b) or a third packet when an error occurs in step b), and d) repeating step c) until the decoding error no longer occurs. The error prevention method has the characteristics of both Type-1 and Type-2 ARQ methods. Therefore, one can obtain constant channel throughput in the burst error containing channel, the random error containing channel, and a channel wherein the two error patterns co-exist simultaneously.



**FOR THE PURPOSES OF INFORMATION ONLY**

Codes used to identify States party to the PCT on the front pages of pamphlets publishing international applications under the PCT.

AL	Albania	ES	Spain	LS	Lesotho	SI	Slovenia
AM	Armenia	FI	Finland	LT	Lithuania	SK	Slovakia
AT	Austria	FR	France	LU	Luxembourg	SN	Senegal
AU	Australia	GA	Gabon	LV	Latvia	SZ	Swaziland
AZ	Azerbaijan	GB	United Kingdom	MC	Monaco	TD	Chad
BA	Bosnia and Herzegovina	GE	Georgia	MD	Republic of Moldova	TG	Togo
BB	Barbados	GH	Ghana	MG	Madagascar	TJ	Tajikistan
BE	Belgium	GN	Guinea	MK	The former Yugoslav Republic of Macedonia	TM	Turkmenistan
BF	Burkina Faso	GR	Greece			TR	Turkey
BG	Bulgaria	HU	Hungary	ML	Mali	TT	Trinidad and Tobago
BJ	Benin	IE	Ireland	MN	Mongolia	UA	Ukraine
BR	Brazil	IL	Israel	MR	Mauritania	UG	Uganda
BY	Belarus	IS	Iceland	MW	Malawi	US	United States of America
CA	Canada	IT	Italy	MX	Mexico	UZ	Uzbekistan
CF	Central African Republic	JP	Japan	NE	Niger	VN	Viet Nam
CG	Congo	KE	Kenya	NL	Netherlands	YU	Yugoslavia
CH	Switzerland	KG	Kyrgyzstan	NO	Norway	ZW	Zimbabwe
CI	Côte d'Ivoire	KP	Democratic People's Republic of Korea	NZ	New Zealand		
CM	Cameroon			PL	Poland		
CN	China	KR	Republic of Korea	PT	Portugal		
CU	Cuba	KZ	Kazakhstan	RO	Romania		
CZ	Czech Republic	LC	Saint Lucia	RU	Russian Federation		
DE	Germany	LI	Liechtenstein	SD	Sudan		
DK	Denmark	LK	Sri Lanka	SE	Sweden		
EE	Estonia	LR	Liberia	SG	Singapore		

## ERROR PREVENTION METHOD FOR MULTIMEDIA

### Technical Field

The present invention relates to an error prevention method for multimedia, and more particularly, to a method for improving data recovery and channel throughput in channels wherein a random error and a burst error occur by using a rate compatible punctured convolutional code (RCPC) and an automatic retransmission on request (ARQ).

### Background Art

Let us consider multimedia terminals which transmit and receive arbitrary packets of data (video, audio, data, or a mixed form of any of those three). The transmitter transmits information packets, for example, I, J, and other packets. For each information packet, the transmitter forms N-sized bit-streams which are different representations of given information packets. For example, the transmitter can generate a packet A (B, C, or D) for given information packet I. Type-1 and Type-2 are different in that they use different retransmission methods. The packets to be transmitted are formed using either convolutional code or RCPC.

FIG. 1 is a block diagram showing a general situation of data transmitted and received using ARQ. The basic concept of Type-1 ARQ will be described as follows with reference to FIG. 1. When a transmitter transmits a packet A having a length of N, a packet decoder 120 in a receiver starts decoding the received packet A 110. At this time, if errors occur in the packet A and no further decoding is possible, for example, channel coding is not employed, channel coding having a 1-bit error or more is employed or, more errors than a channel coder can detect and correct occur, the receiver asks the transmitter to send the same packet A again. Here, retransmission would be repeated either until the decoder 120 receives an error-free packet A, or for some specific number of iterations to perform transmission and receiving with respect to the next packet. Type-1 ARQ is very effective in burst-error containing channel. Next, Type-2 ARQ will be described. Up to now, there are three types of Type-2 ARQ, i.e., a basic type, a Class A and a Class B, each of which uses RCPC given information I (J, K, ...).

FIG. 2 is a conceptual diagram showing the operation of the basic type, wherein arrows represent combination. Here, given information I, the transmitter generates packets A and B using RCPC at a rate of  $\frac{1}{2}$  and transmits only the packet A. The decoder in the receiver attempts to decode the packet A. If successful, the decoder then attempts to decode the first packet of two for the next information J. Otherwise, the receiver asks the transmitter to send the packet B. Also, the decoder attempts decoding a combination of packets A and B. If successful, the decoder then attempts to decode the first packet of two for the next information J. Otherwise, the receiver asks the transmitter to send the packet A again and all of these procedures are repeated. The basic type has an advantage in that implementation is not so complicated.

FIG. 3 is a conceptual diagram showing the operation of the Class A packet (Lin-Yu), wherein \* denotes self-decoding and arrows represent combination. The operational principle thereof is similar to the basic type except how to combine packets A and B when both packets fail to be decoded. That is, the decoder attempts decoding the combination of packets A and B, and if it fails, the receiver asks the transmitter to send the packet A again. Next, if the decoder succeeds decoding only the packet A, the next information J is processed, and if the decoder fails, the receiver combines the previously stored packet B and currently received packet A (i.e., in general, interleaves the two) to attempt decoding. This method is more effective in a random error containing channel rather than in a burst error containing channel.

Next, Class B is significantly more complicated than the basic type and the Class A. The basic concept thereof is based on the Class A. First, the Class A (Lin-Yu) is performed by generating the packets A and B given the information I using RCPC at a rate of  $\frac{1}{2}$ . As described above, the Type-1 ARQ is greatly effective in the burst error containing channel. However, with Type-1 ARQ, retransmission would be more frequent in the random error containing channel, which causes drastically lower channel throughput. Even though Type-2 ARQ allows good performance in the random error containing channel, retransmission would be more frequent in the burst error containing channel: therefore, channel throughput can be lowered.

### Disclosure of Invention

It is an object of the present invention to provide a method for maintaining channel throughput at a certain level in a random error containing channel and a burst error containing channel by operating like Type-1 in the burst error containing channel, while operating similarly to a basic type or a Class A of Type-2 in the random error containing channel.

To accomplish the above object, there is provided an error prevention method in a method for decoding a plurality of packets of given information, comprising the steps of a) decoding one of the plurality of packets, b) decoding another packet when an error occurs in the decoding in step a), c) decoding a combination of the decoding error packets when an error occurs in step b) or the third packet, and d) repeating step c) until the decoding error no longer occurs.

### Brief Description of the Drawings

The above object and advantage of the present invention will become more apparent by describing in detail a preferred embodiment thereof with reference to the attached drawings in which:

FIG. 1 is a block diagram showing the general situation of data transmission and reception using an ARQ method;

FIG. 2 is a conceptual diagram showing the operation of a basic type;

FIG. 3 is a conceptual diagram showing the operation of a Class A;

FIG. 4 is a block diagram showing the structure of an apparatus for realizing error prevention according to the present invention;

FIG. 5 is a conceptual diagram showing a procedure for processing of received packets A, B, C and D in a decoder of a receiver shown in FIG. 4; and

FIG. 6 is a flow chart showing a procedure for processing received packets in a decoder according to the present invention.

### Best Mode For Carrying out the Invention

The present invention is a method for performing a hybrid-type ARQ which mixes Type-1 and Type-2 methods.

Referring to FIG. 4, an error prevention apparatus includes: a transmitter provided with a packet buffer 430 for producing packets A, B, C and D using an

RCPC 420 having a rate of  $1/4$  for given information packet I 410; an inverse RCPC 440; and a receiver provided with a buffer 450 for storing the received packet, for sending an ARQ and a packet number to the transmitter via a channel. In FIG. 4, an RCPC logic is fixed at a rate of  $1/4$ . A portion comprised of four arbitrary polynomials which satisfies a locally invertible characteristic simultaneously produces RCPC-processed packets A, B, C and D. Also, the transmitter has a maximum of four retransmissions. Here, the local inversion in the RCPC denotes that original information I can be obtained with any one of the packets A and B and with a combination of the packets A and B.

FIG. 5 is a conceptual view showing a procedure for processing the received packets A, B, C and D in the decoder of the receiver shown in FIG. 4, wherein \* indicates self-decoding and a bracket indicates a combination of packets (generally, an interleaving operation).

FIG. 6 is a flow chart outlining a procedure for processing the received packets in a decoder according to the present invention.

As shown in FIG. 6, the transmitter produces packets A, B, C and D using the RCPC 420 in step 612. The first packet A is transmitted to the receiver in step 614. The decoder attempts decoding the packet A in step 616. In step 616, if the packet A is decoded, the decoded results are stored in the buffer 450 (step 642) and the job for other information (e.g., information J) proceeds (step 644); otherwise, an ARQ signal is sent to the transmitter to request transmission of the packet B (step 618). The decoder attempts decoding only the packet B in step 620. If successful, the decoded results are stored in the buffer 450 in step 642 and the job for other information (e.g., information J) proceeds in step 644. If the decoder fails, it then attempts to decode the combination of the packets A and B in step 622, which is indicated by \*AB in FIG. 5. At this time, if the combination of the packets A and B as shown in FIG. 5 are decoded, the decoded results are stored in the buffer 450 (step 642) and the job for other information (e.g., information J) proceeds (step 644). If the combination of the packets A and B is not decoded, the transmitter is requested to transmit the packet C by sending the ARQ signal thereto in step 624. The decoder then attempts decoding only packet C in step 626. If successful, the decoded results are stored in the buffer 450 in step 642 and the job for other information (e.g., information J) proceeds in step 644. Otherwise, the

decoder attempts to decode the combination of the packets B and C in step 628, which is indicated by \*BC in FIG. 5. If the combination of the packets B and C is successfully decoded, the decoded results are stored in the buffer 450 (step 642) and the job for other information (e.g., information J) proceeds (step 644). If the combination thereof is not decoded, the combination of packets B and C is combined with packet A as shown in FIG. 5 as \*ABC and decoding is attempted in step 630. Here, if the combination of packets A, B and C is decoded, the decoded results are stored in the buffer 450 (step 642) and the job for other information (e.g., information J) proceeds (step 644). Otherwise, the transmitter is requested to transmit the packet D by sending the ARQ signal thereto in step 632. The decoder then attempts decoding only the packet D in step 634. If successful, the decoded results are stored in the buffer 450 in step 642 and the job for other information (e.g., information J) proceeds in step 644. Otherwise, the decoder attempts to decode the combination of packets C and D in step 636, which is indicated by \*CD in FIG. 5. If the combination of packets C and D is decoded, the receiver stores the decoded results in the buffer 450 (step 642) and performs the job for other information (e.g., information J) (step 644). Otherwise, the receiver combines packets C and D with packet B as indicated by \*BCD in FIG. 5 and attempts decoding the combination in step 638. Here, if the combination of packets B, C and D is decoded, the decoded results are stored in the buffer 450 (step 642) and the job for other information (e.g., information J) proceeds (step 644). Otherwise, the receiver combines packets B, C and D with packet A as indicated by \*ABCD in FIG. 5 and attempts decoding combination in step 640. If the combination of packets A, B, C and D is decoded, the receiver stores the decoded results in the buffer 450 (step 642) and performs the job for other information (e.g., information J) (step 644). Otherwise, the process returns to step 614 to repeat all of these procedures until no errors occur. Meanwhile, the receiver stores the decoded results in the buffer 450 in step 642 and performs the job for the next information (e.g., information J, K....) in step 644.

#### Industrial Applicability

As described above, the present invention has the characteristics of both Type-1 and Type-2 ARQ methods: therefore, one can obtain constant channel

throughput in the burst error containing channel, the random error containing channel, and a channel where the two error patterns coexist simultaneously. In the burst error containing channel, the method of the present invention is performed nearly the same as or better than the Type-1 method and much better than the Type-2 method. As for the random error containing channel, since the method of the present invention is performed similar to the Type-2 method, it also performs almost the same as the Type-2 method, but much better than the Type-1 method.



What is claimed is:

1. A method for preventing errors during the decoding a plurality of packets of given information, comprising the steps of:

a) decoding one of said plurality of packets;

b) decoding another packet when an error occurs during said decoding in said step a);

c) decoding a combination of said packets of said steps a) and b) or a third packet when an error occurs during said decoding in said step b); and

d) repeating said step c) until said decoding error no longer occurs.

2. An error prevention method as claimed in claim 1, further comprising the step of, storing said decoded results and standing by decoding of a plurality of packets of the next information when said decoding error no longer occurs during said steps a) to d).

3. An error protecting method as claimed in claim 1, wherein, at least two decoding error packets are combined and decoded when the number of said packets in said step c) is at least three.

4. An error prevention method as claimed in claim 3, wherein decoding is sequentially performed starting from when the number of combined packets is two and including a latest packet determined to be a decoding error packet.

1/3

FIG. 1 (PRIOR ART)

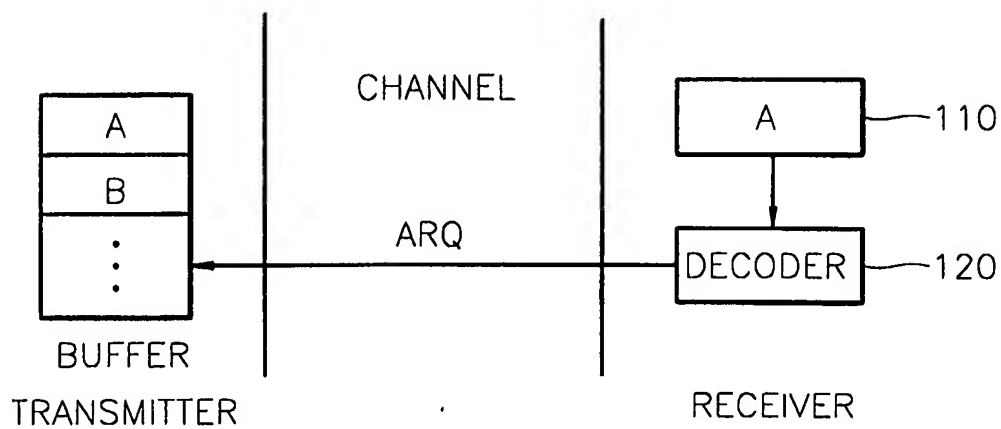


FIG. 2 (PRIOR ART)

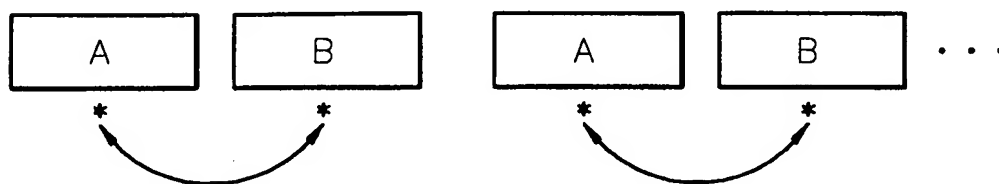
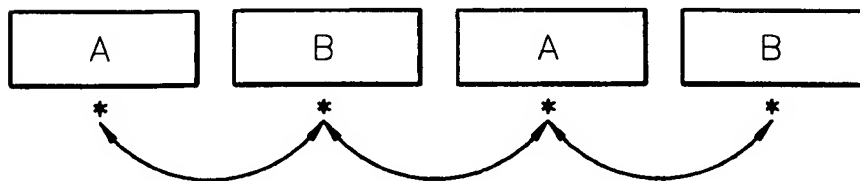


FIG. 3 (PRIOR ART)



2/3

FIG. 4

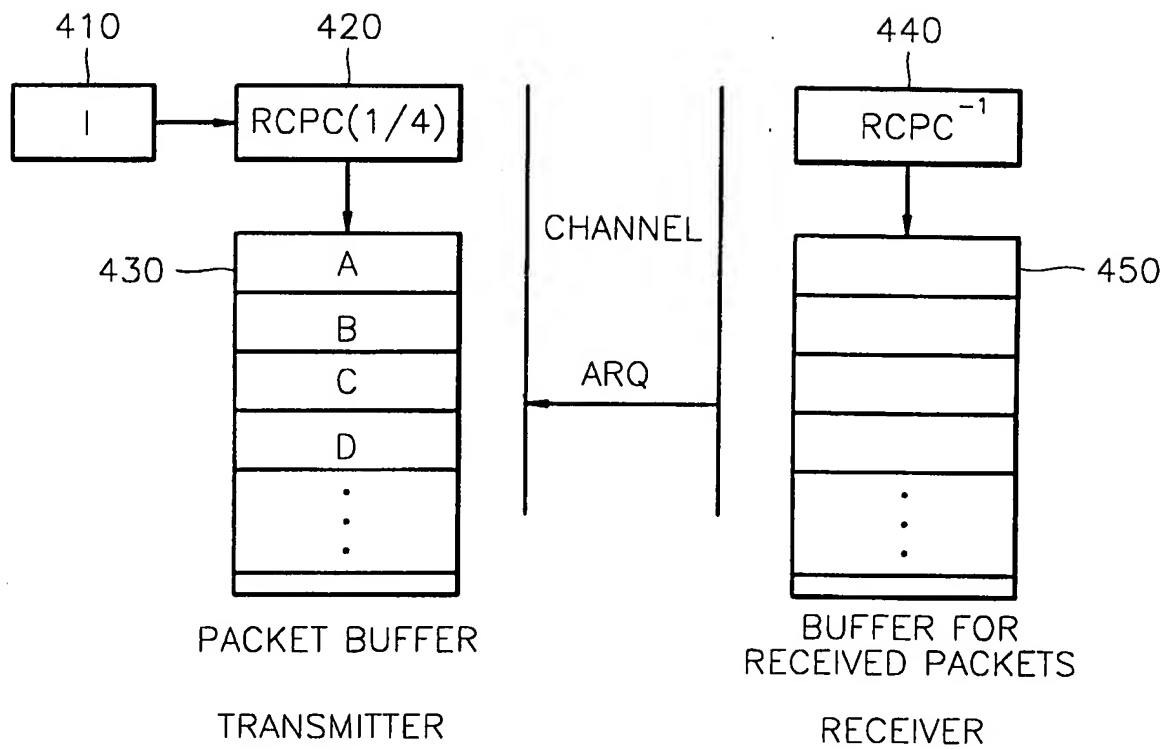
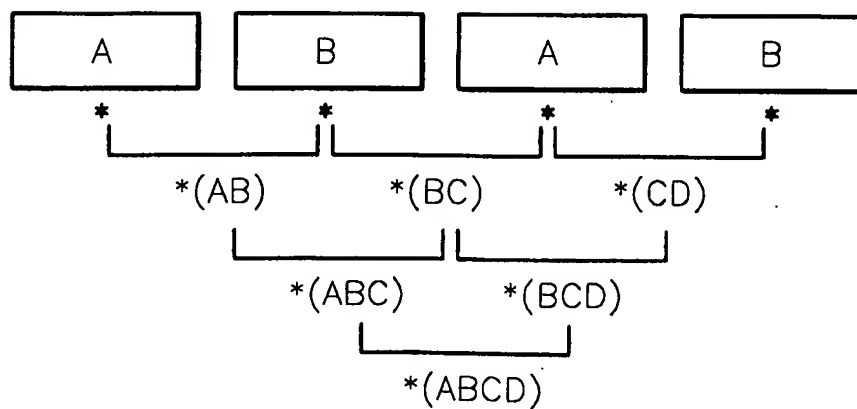
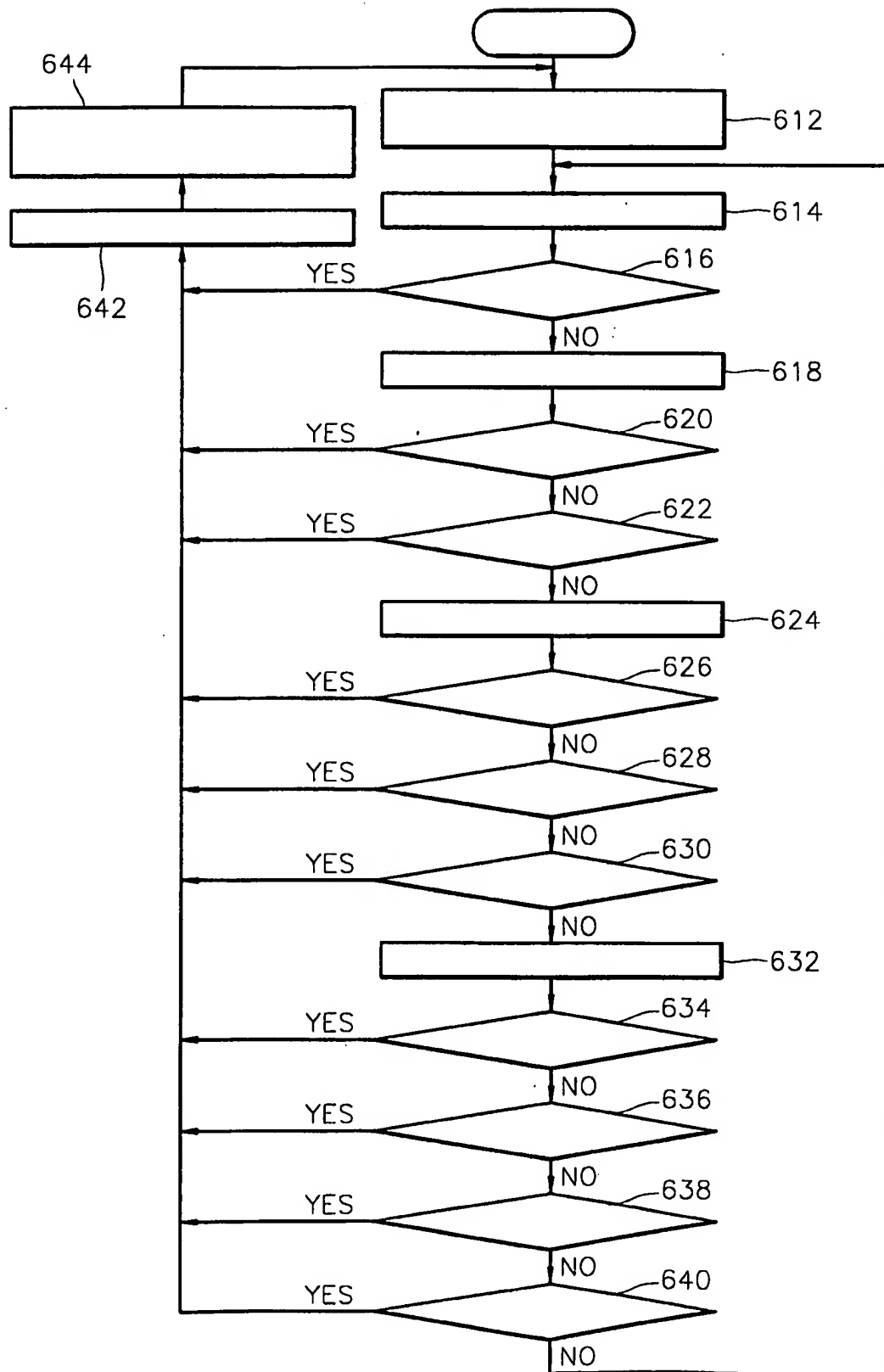


FIG. 5



3/3

FIG. 6



## INTERNATIONAL SEARCH REPORT

International application No.

PCT/KR 98/00004

## A. CLASSIFICATION OF SUBJECT MATTER

IPC<sup>6</sup>: H 03 M 13/00

According to International Patent Classification (IPC) or to both national classification and IPC

## B. FIELDS SEARCHED

Minimum documentation searched (classification system followed by classification symbols)

IPC<sup>6</sup>: H 03 M 13/00, 13/02, 13/12, 13/22; H 04 L 1/00, 1/12, 1/14, 1/16

Documentation searched other than minimum documentation to the extent that such documents are included in the fields searched

Electronic data base consulted during the international search (name of data base and, where practicable, search terms used)

WPIL

## C. DOCUMENTS CONSIDERED TO BE RELEVANT

Category*	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
A	US 5 487 068 A (JEFFREY C. SMOLINSKA et al.) 23 January 1996 (23.01.96), totality. -----	1

☐ Further documents are listed in the continuation of Box C.☒ See patent family annex.

## \* Special categories of cited documents:

"A" document defining the general state of the art which is not considered to be of particular relevance

"E" earlier document but published on or after the international filing date

"L" document which may throw doubts on priority claim(s) or which is cited to establish the publication date of another citation or other special reason (as specified)

"O" document referring to an oral disclosure, use, exhibition or other means

"P" document published prior to the international filing date but later than the priority date claimed

"T" later document published after the international filing date or priority date and not in conflict with the application but cited to understand the principle or theory underlying the invention

"X" document of particular relevance; the claimed invention cannot be considered novel or cannot be considered to involve an inventive step when the document is taken alone

"Y" document of particular relevance; the claimed invention cannot be considered to involve an inventive step when the document is combined with one or more other such documents, such combination being obvious to a person skilled in the art

"&amp;" document member of the same patent family

Date of the actual completion of the international search

03 April 1998 (03.04.98)

Date of mailing of the international search report

28 April 1998 (28.04.98)

Name and mailing address of the ISA/ AT

AUSTRIAN PATENT OFFICE

Kohlmarkt 8-10

A-1014 Vienna

Facsimile No. 1/53424/535

Authorized officer

Zugarek

Telephone No. 1/53424/383

**INTERNATIONAL SEARCH REPORT**  
Information on patent family members

International application No.

PCT/KR 98/00004

Im Recherchenbericht angeführtes Patentdokument Patent document cited in search report Document de brevet cité dans le rapport de recherche	Datum der Veröffentlichung Publication date Date de publication	Mitglied(er) der Patentfamilie Patent family member(s) Membre(s) de la famille de brevets	Datum der Veröffentlichung Publication date Date de publication
US A 5487068	23-01-96	AU A1 29011/95	04-03-96
		AU B2 675038	16-01-97
		CA AA 2171015	15-02-96
		CN A 1131488	18-09-96
		EP A1 727119	21-08-96
		EP A4 727119	20-08-97
		WO A1 9604736	15-02-96